

# An Introduction to Inductive Inference

Nan Fang

Institut für Informatik, Universität Heidelberg

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# An introductory example

Learner reads more and more data about an r.e. set and outputs hypotheses what the set is.

Data	Hypotheses
2	Set of even numbers;
2,3	Set of all numbers;
2,3,5	Set of prime numbers;
2,3,5,13	Set of prime numbers;
2,3,5,13,1	Set of Fibonacci numbers;
2,3,5,13,1,8	Set of Fibonacci numbers;
2,3,5,13,1,8,21	Set of Fibonacci numbers.

Successful learning means that the conjectures converge to one correct concept.

Think about language acquisition for children:

- Children can master any natural language in a few years' time on the basis of rather casual and unsystematic exposure to it.
- Many linguists (notably, Chomsky[1975],[1986]) believe that children are not genetically prepared to acquire any arbitrary language.
- Instead, a relatively small class  $\mathcal{H}$  of languages may be singled out as “humanly possible”.
- While linguists fail to propose a nontrivial description of  $\mathcal{H}$ .
- The branch of linguistics known as “comparative grammar” is the attempt to characterize the class of (biologically possible) natural languages through formal specification of their grammars.

# A further example

- Let  $\mathcal{L} = \{\omega \setminus \{n\} : n \in \omega\}$ .
- Find an effective process to learn (identify) every member of  $\mathcal{L}$ .
  
- How about the class  $\mathcal{L} \cup \{\omega\}$ ?

# A learning paradigm

A formal model of inductive inference should contain the following concepts:

- 1 a theoretically possible reality
- 2 intelligible hypotheses
- 3 the data available about any given reality, were it actual
- 4 a learner
- 5 successful behaviour by a learner working in a given, possible reality

# Our models

- 1 A theoretically possible reality is a family of languages coded by a class of computably enumerable sets.
- 2 Hypotheses are symbolic representations of a real or fictitious world, for our reality we take a Gödel numbering.
- 3 The information about a given reality can be of different type, mainly two. One consist of only positive instances, the other contains negative ones as well.
- 4 A learner is a Turing machine converting the data he observed to hypotheses. Usually, we require it to be a total Turing machine. We could also have some other requirement on him.
- 5 One possible successful behaviour could be that, after many attempts, the learner's hypotheses finally converges to a code of the real object.

## Definition

- we call a computably enumerable subset of  $\omega$  a *language*.
- A *text* is a total function  $t : \omega \rightarrow \omega \cup \{\#\}$ .
- $content(t) = range(t) \setminus \{\#\}$ .
- A text  $t$  is a *text for the set*  $A \subseteq \omega$  if  $A = content(t)$ .
- $SEQ = (\omega \cup \{\#\})^*$
- A *learner* is a total computable function  $M : SEQ \rightarrow \omega \cup \{?\}$ .

## Definition

- The learner  $M$  *TxtEx-learns* the language  $L$  from the text  $t$  for  $L$  if there is a number  $n_0$  such that

$$\forall n \geq n_0 (M(t \upharpoonright n) = M(t \upharpoonright n_0) \in \omega) \& L = W_{M(t \upharpoonright n_0)}$$

- The learner  $M$  *TxtEx-learns* the language  $L$  if  $M$  *TxtEx-learns*  $L$  from all texts  $t$  for  $L$ .
- The learner  $M$  *TxtEx-learns* the family  $\mathcal{L}$  of languages if  $M$  *TxtEx-learns* all languages in  $\mathcal{L}$ .
- A family  $\mathcal{L}$  of languages is *TxtEx-learnable* if there is a learner  $M$  which *TxtEx-learns* the family  $\mathcal{L}$ . The class of all *TxtEx-learnable* families is denoted by *TxtEx*.



# Some properties of TxtEx

## Proposition (Monotonicity)

*If  $\mathcal{L}_0 \subset \mathcal{L}_1$  and  $\mathcal{L}_1 \in \text{TxtEx}$ , then  $\mathcal{L}_0 \in \text{TxtEx}$ .*

## Proposition

*Any finite family  $\mathcal{L}$  of languages is in TxtEx.*

## Proposition

*The family  $\mathcal{FIN}$  of all finite sets is in TxtEx.*

# Other learning criteria

## Definition (finite learning)

The learner  $M$  *TxtFin-learns* the language  $L$  from the text  $t$  for  $L$  if there is a number  $n_0$  such that

$$\forall n \geq n_0 (M(t \upharpoonright n) = M(t \upharpoonright n_0) \in \omega)$$

$$\& \forall n < n_0 (M(t \upharpoonright n) = ?) \& L = W_{M(t \upharpoonright n_0)}.$$

## Definition (behaviourally correct learning)

The learner  $M$  *TxtBC-learns* the language  $L$  from the text  $t$  for  $L$  if there is a number  $n_0$  such that

$$\forall n \geq n_0 (M(t \upharpoonright n) \in \omega \& L = W_{M(t \upharpoonright n_0)}).$$

Just as in the case of TxtEx-learning, we define the classes TxtFin and TxtBC accordingly.

# Properties of learning from text

## Proposition (Monotonicity)

*Let  $X \in \{\text{TxFin}, \text{TxEEx}, \text{TxBC}\}$ . If  $\mathcal{L}_0 \subset \mathcal{L}_1$  and  $\mathcal{L}_1 \in X$ , then  $\mathcal{L}_0 \in X$ .*

## Proposition

$$\text{TxFin} \subseteq \text{TxEEx} \subseteq \text{TxBC}$$

## Proposition

*Let  $L_0$  and  $L_1$  be languages such that  $L_0 \subset L_1$ . Then  $\{L_0, L_1\} \notin \text{TxFin}$*

## Definition

- An *informant* for a language  $L$  is its characteristic sequence  $L(0)L(1)L(2)\dots$
- The learner  $M$  *InfEx-learns* the language  $L$  if there is a number  $n_0$  such that

$$\forall n \geq n_0 (M(L \upharpoonright n) = M(L \upharpoonright n_0) \in \omega) \& L = W_{M(L \upharpoonright n_0)}$$

- The learner  $M$  *InfEx-learns* the family  $\mathcal{L}$  of languages if  $M$  InfEx-learns all languages in  $\mathcal{L}$ .
- A family  $\mathcal{L}$  of languages is *InfEx-learnable* if there is a learner  $M$  which InfEx-learns the family  $\mathcal{L}$ . The class of all InfEx-learnable families is denoted by InfEx.

Similarly, we can define the classes InfFin and InfBC.

## Proposition

$$\text{InfFin} \subseteq \text{InfEx} \subseteq \text{InfBC}$$

## Proposition

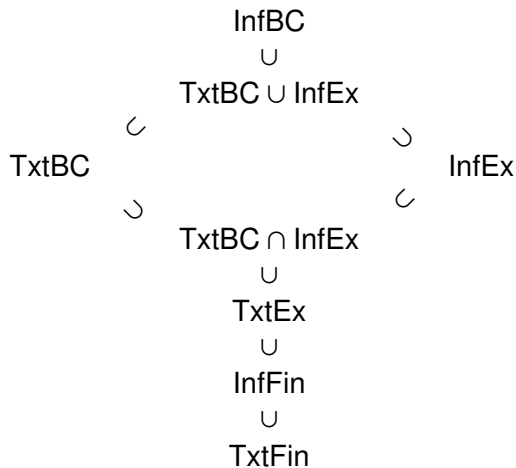
*For  $X = \text{Ex}, \text{Fin}, \text{BC}$ ,*

$$\text{TxtX} \subseteq \text{InfX}.$$

Questions: What's the relationship among these six major classes? Are their inclusion relations strict?

# Learnability Hierarchy

The answer is yes, and there is a full hierarchy of them:



# Other variants of learning model

- A learner may be not totally computable; learner(machine) may have some time or space requirement.
- The hypotheses space is not a Gödel numbering.
- Oracle learning
- Probabilistic learning

# Computational Learning VS. Statistical Learning

Computational Learning

VS.

Statistical Learning

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*What can machines learn, exactly, in theory?*

*How can machines learn, correctly (less error and more efficiently), in practice?*

Linguist

Language teacher



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