An Introduction to Forcing

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- Preliminary
- Intention

Proving and Consistency Proofs

- Generic Extension
- Forcing Relation
- Forcing with Finite Partial Functions

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Preliminary

Our Logic

• Soundness and Completeness of first-order predicate logic *T* is consistent if and only if *T* has a (countable) model.

Preliminary

Gödel's incompleteness results
 We can only hope relative consistency results, e.g.

 $Con(ZFC) \rightarrow Con(ZFC + V = L)$

- Our Theory of Sets
 - Axioms of ZFC
 - Partial orders, boolean algebras, filters, dense sets, chain/antichain, etc.
 - Relativization and absoluteness
 - Others: Δ-system, cardinal arithmetic, etc.

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• We have found a "model" L of constructible sets in the ground model and shown that

$$(ZF + V = L)^{L}, V = L \rightarrow GCH \wedge AC.$$
 (1)

- No inner model can be found to make (ZF + V ≠ L) true in it as long as ZFC is consistent.
- We should extend our ground model *M* to be *M*[*G*], the generic extension.

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Basic Idea

- Start from *M*, a countable, transitive model of ZFC
- Design a partial order \mathbb{P} (the set of conditions) in M
- Pick a generic filter $G \subseteq \mathbb{P}$, usually $G \notin M$.
- Make *M*[*G*] the smallest transitive model of ZFC containing both *M* and *G*.
- The truth in M[G] base mainly on the ground model M and the partial order \mathbb{P} .

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Foundational Theorem of Forcing

Theorem (Theorem of Generic Model)

Given ground model M, partial order $\mathbb{P} \in M$, and generic filter G, there is a M[G] such that

- *M*[*G*] is a transitive model of ZFC;
- *M* ⊆ *M*[*G*] and *G* ∈ *M*[*G*];
- *M*[*G*] is the smallest such model.

Theorem (Forcing Theorem)

Under the hypotheses of the previous theorem. Given formula $\varphi(v_1, \ldots, v_n)$ and \mathbb{P} -name $\tau_1, \ldots, \tau_n \in M$.

 $\varphi(\tau_1^G,\ldots,\tau_n^G)^{M[G]}$ if and only if $\exists p \in G(p \Vdash \varphi(\tau_1,\ldots,\tau_n))^M$.

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Generic Filter

Definition

 $G \subseteq \mathbb{P}$ is a generic filter if *G* is a filter and for each dense $D \subseteq \mathbb{P}$ such that $D \in M$, $G \cap D \neq \emptyset$.

- We can always found a generic filter in a countable ground model.
- We can do forcing from arbitrary partial order ℙ, but only the following case is nontrivial.

For each $p \in \mathbb{P}$, there are $q \leq p$ and $r \leq p$ such that $q \perp r$.

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ℙ-name

People in *M* should think about possible extensions, and denote the objects in them by \mathbb{P} -names.

Definition

 τ is \mathbb{P} -name if τ is a relation, and for all $(\pi, p) \in \tau$, π is \mathbb{P} -name, $p \in \mathbb{P}$.

- The definition of $\mathbb P$ must be considered as inductive.
- \mathbb{P} -name is an absolute notion.
- \mathbb{P} -names can be ranked.

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The Generic Extension *M*[*G*]

We define the object τ^{G} that the name τ denotes, and *G* assigns to.

Definition

 τ is \mathbb{P} -name,

$$au^{m{G}} = ig\{\pi^{m{G}} \mid (\exists m{p} \in m{G})(\pi,m{p}) \in auig\}.$$

Note that the definition is also inductive. The generic extension is defined as,

Definition

$$\boldsymbol{M}[\boldsymbol{G}] = \big\{ \boldsymbol{\tau}^{\boldsymbol{G}} \mid \boldsymbol{\tau} \in \boldsymbol{M}^{\mathbb{P}} \big\}.$$

M[G] is transitive if M is.

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The Canonical Names

Definition

For each set x in the ground model, we define

$$\check{x} = \big\{ (\check{y}, \boldsymbol{\rho}) \mid \boldsymbol{y} \in \boldsymbol{x}, \boldsymbol{\rho} \in \mathbb{P} \big\}.$$
(4)

We claim that $\check{x}^G = x$. Thus $M \subseteq M[G]$.

Definition

$$\dot{G} = \{(\check{p}, p) \mid p \in \mathbb{P}\}.$$

(5)

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- G has nothing to do with G.
- $G = \dot{G}^G \in M[G].$
- G is the oracle beyond M and finally decides M[G].

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Forcing Relation

We define the forcing relation $p \Vdash \varphi(\tau_1, \ldots, \tau_n)$, where p is a condition, τ_1, \ldots, τ_n are \mathbb{P} -name.

Definition

For atomic formula, we define by induction on (rank(\(\tau_1\), rank(\(\tau_2\)))

•
$$p \Vdash \tau_1 = \tau_2$$
 iff $p \Vdash \tau_1 \subseteq \tau_2$ and $p \Vdash \tau_2 \subseteq \tau_1$,
 $p \Vdash \tau_1 \subseteq \tau_2$ iff for each $(\pi, r) \in \tau_1$,
 $\{q \mid q \leq r \rightarrow q \Vdash \pi \in \tau_1\}$ is dense below p ;

- $p \Vdash \tau_1 \in \tau_2$ iff $\{q \mid \exists (\pi, r) \in \tau_2 (q \leq r \land q \Vdash \pi = \tau_1)\}$ is dense below p.
- The atomic case is absolute.

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Forcing Relation Boolean and Quantifier Cases

We continue the definition by induction on the complexity of formula.

Definition

- ◎ $p \Vdash \neg \varphi$ if and only if for each $q \leq p$, $q \nvDash \varphi$;

• $p \Vdash \exists x \varphi(x)$ if and only if $\{q \mid \exists \pi(\pi \text{ is } \mathbb{P}\text{-name } \land q \Vdash \varphi(\pi))\}$ is dense below p.

- Forcing relation is not absolute generally.
- The forcing relation is the "logic" of the people living in *M*. It decides the outline of every possible *M*[*G*].

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Some Additional Property of Forcing Relation

Theorem

- If $q \leq p$, then $p \Vdash \varphi$ implies $q \Vdash \varphi$.
- $\{p \mid p \Vdash \varphi \lor p \Vdash \neg \varphi\}$ is dense.
- No $p \in \mathbb{P}$ forces both φ and $\neg \varphi$.

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Proof of the Forcing Theorem

We prove that

$$\varphi(\tau_1^G,\ldots,\tau_n^G)$$
 iff $\exists p \in G(p \Vdash \varphi(\tau_1,\ldots,\tau_n))^M$

- Atomic case
- Boolean and quantifier case

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Finish the Proof of the Generic Model Theorem

Lemma

 $M[G] \vDash ZFC.$

Proof.

- Extensionality: *M*[*G*] is transitive.
- Foundation: holds in each \in model.
- For those axioms that asserts existence of sets, we should design appropriate names.

Lemma

If N is a transitive model of ZFC and that $M \subseteq N$, $G \in N$, then $M[G] \subseteq N$.

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Finite Partial Functions

The crucial trick is to design the partial order. Here we give a simple example.

Definition (Finite partial functions)

$$Fn(I,J) = \{p: | |p| < \omega \land p \text{ is a function } \land \text{dom } p \subseteq I \land \text{ran } p \subseteq J\}.$$
(6)

The order on Fn(I, J) is defined as

$$p \leq q \leftrightarrow p \supseteq q.$$
 (7)

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Some Example

• Forcing with $Fn(\omega, \omega_1)$

 $\bigcup G \text{ is a total function mapping } \omega \text{ onto } \omega_1?$ (8)

- Forcing with $Fn(\kappa \times \omega, 2)$
 - $f = \bigcup G : \kappa \times \omega \mapsto 2$ is total.
 - For $\alpha < \kappa$, Letting

 $f_{\alpha}: \omega \mapsto 2$, such that $f_{\alpha}(n) = f(\alpha, n)$. (9)

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⟨f_α : α < κ⟩ is an one-to-one sequence mapping κ into 2^ω in *M*[*G*].

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Preserving Cardinals

Theorem

 $\mathbb{P} \in M$. If $(\mathbb{P} \text{ is c.c.c.})^M$, then for each generic G of \mathbb{P} over M and ordinal $\alpha \in M$,

 $(\alpha \text{ is a cardinal})^M \leftrightarrow (\alpha \text{ is a cardinal})^{M[G]}.$

Lemma

 $Fn(\kappa \times \omega, 2)$ is c.c.c..

Use the Δ -system theorem.

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Further Discussion

Theorem

For each κ with cf $\kappa > \omega$.

$$\operatorname{Con}(ZFC) \to \operatorname{Con}(ZFC + 2^{\omega} = \kappa).$$

- We have shown that Con(ZFC) → Con(ZFC + 2^ω > κ) for each κ.
- Forcing with a ground model of ZFC + GCH.

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Summary

- The generic extension M[G] is built from the ground model M, the partial order $\mathbb{P} \in M$, and the generic filter G (usually not in M).
- The The general truth in M[G] is already described by the forcing relation in M, so is decided mainly by \mathbb{P} and M.
- Outlook
 - Proper Forcing.
 - *P_{max}* Forcing.

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For Further Reading





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