

An introduction to forcing

Wang Yunsong

School of Mathematical Science

1600010615@pku.edu.cn

September 17, 2019

Overview

1 Generic extensions

2 Forcing

Our goal: $Con(ZFC) \rightarrow Con(ZFC + \neg CH)$

Our goal: $Con(ZFC) \rightarrow Con(ZFC + \neg CH)$

Theorem(Godel)

$$T \vdash \varphi \Leftrightarrow T \models \varphi$$

Our goal: $Con(ZFC) \rightarrow Con(ZFC + \neg CH)$

Theorem(Godel)

$$T \vdash \varphi \Leftrightarrow T \models \varphi$$

So if there is a $M \models ZFC$, we want to find a new model $N \models ZFC + \neg CH$

We will construct a model $M[G]$, s.t $G \notin M$, $M \subseteq M[G]$

– M transitive set, $x \in M \wedge y \in x \Rightarrow y \in M$

- M transitive set, $x \in M \wedge y \in x \Rightarrow y \in M$
- A p.o. of M is a triple, $\langle \mathbb{P}, \leq, \mathbb{1} \rangle$ s.t. \leq partially orders \mathbb{P} and $\mathbb{1}$ is a largest element of \mathbb{P} (i.e., $\forall p \in \mathbb{P}(p \leq \mathbb{1})$), and $\mathbb{P} \in M$.

- M transitive set, $x \in M \wedge y \in x \Rightarrow y \in M$
- A p.o. of M is a triple, $\langle \mathbb{P}, \leq, \mathbb{1} \rangle$ s.t. \leq partially orders \mathbb{P} and $\mathbb{1}$ is a largest element of \mathbb{P} (i.e., $\forall p \in \mathbb{P}(p \leq \mathbb{1})$), and $\mathbb{P} \in M$.
- Let $G \subseteq \mathbb{P}$, we call G is a filter on \mathbb{P} :
 - $G \neq \emptyset$
 - $\forall p, q \in \mathbb{P}((p \in G \wedge q \in G) \rightarrow \exists r \in G(r \leq p \wedge r \leq q))$
 - $\forall p, q \in \mathbb{P}((p \leq q \wedge p \in G) \rightarrow q \in G)$

- $D \subseteq \mathbb{P}$, D is dense in \mathbb{P} , if $\forall p \in \mathbb{P}, \exists q \in D (q \leq p)$.

- $D \subseteq \mathbb{P}$, D is dense in \mathbb{P} , if $\forall p \in \mathbb{P}, \exists q \in D (q \leq p)$.
- G is \mathbb{P} - *generic* over M iff G is a filter on \mathbb{P} and for all dense $D \subseteq \mathbb{P}, D \in M \rightarrow G \cap D \neq \emptyset$.

- $D \subseteq \mathbb{P}$, D is dense in \mathbb{P} , if $\forall p \in \mathbb{P}, \exists q \in D (q \leq p)$.
- G is \mathbb{P} – generic over M iff G is a filter on \mathbb{P} and for all dense $D \subseteq \mathbb{P}, D \in M \rightarrow G \cap D \neq \emptyset$.

Lemma

if M is countable and $p \in \mathbb{P}$, then there is a G which is \mathbb{P} – generic over M such that $p \in G$.

- $D \subseteq \mathbb{P}$, D is dense in \mathbb{P} , if $\forall p \in \mathbb{P}, \exists q \in D (q \leq p)$.
- G is \mathbb{P} – generic over M iff G is a filter on \mathbb{P} and for all dense $D \subseteq \mathbb{P}, D \in M \rightarrow G \cap D \neq \emptyset$.

Lemma

if M is countable and $p \in \mathbb{P}$, then there is a G which is \mathbb{P} – generic over M such that $p \in G$.

Proof: M is countable, so let $D_n (n \in \omega)$ enumerate all dense subsets of \mathbb{P} which are in M .

Then choose a sequence $q_n (n \in \omega)$ so that $p = q_0 \geq q_1 \geq \dots$ and $q_{n+1} \in D_n$. Let G be the filter generated by $\{q_n : n \in \omega\}$.

- $D \subseteq \mathbb{P}$, D is dense in \mathbb{P} , if $\forall p \in \mathbb{P}, \exists q \in D (q \leq p)$.
- G is \mathbb{P} – generic over M iff G is a filter on \mathbb{P} and for all dense $D \subseteq \mathbb{P}, D \in M \rightarrow G \cap D \neq \emptyset$.

Lemma

if M is countable and $p \in \mathbb{P}$, then there is a G which is \mathbb{P} – generic over M such that $p \in G$.

Proof: M is countable, so let $D_n (n \in \omega)$ enumerate all dense subsets of \mathbb{P} which are in M .

Then choose a sequence $q_n (n \in \omega)$ so that $p = q_0 \geq q_1 \geq \dots$ and $q_{n+1} \in D_n$. Let G be the filter generated by $\{q_n : n \in \omega\}$.

Remark: Dense subset can be replaced by open dense subset or maximum anti-chain.

Lemma

If M is a transitive model of ZFC, $\mathbb{P} \in M$ is a p.o. such that $\forall p \in \mathbb{P}$
 $\exists q, r \in \mathbb{P} (q \leq p \wedge r \leq p \wedge q \perp r)$, and G is \mathbb{P} -generic over M , then $G \notin M$.

Lemma

If M is a transitive model of ZFC, $\mathbb{P} \in M$ is a p.o. such that $\forall p \in \mathbb{P} \exists q, r \in \mathbb{P} (q \leq p \wedge r \leq p \wedge q \perp r)$, and G is \mathbb{P} -generic over M , then $G \notin M$.

Remark: If condition fails for \mathbb{P} , then there is a filter G on \mathbb{P} which intersects all dense subsets of \mathbb{P} , and if $\mathbb{P} \in M$, then $G \in M$.

Then forcing to such a \mathbb{P} will be trivial. Thus, almost all p.o. we considered satisfy this condition, although it is never needed in the abstract treatment of forcing.

Let M be a c.t.m for ZFC, with \mathbb{P} a p.o. in M and G is \mathbb{P} – *generic* over M .

Let M be a c.t.m for ZFC, with \mathbb{P} a p.o. in M and G is \mathbb{P} – *generic* over M .

We shall show how to construct another c.t.m for ZFC, called $M[G]$, which will satisfy $M \subset M[G]$, $\alpha(M) = \alpha(M[G])$, and $G \in M[G]$.

Let M be a c.t.m for ZFC, with \mathbb{P} a p.o. in M and G is \mathbb{P} – *generic* over M .

We shall show how to construct another c.t.m for ZFC, called $M[G]$, which will satisfy $M \subset M[G]$, $\alpha(M) = \alpha(M[G])$, and $G \in M[G]$.

$M[G]$ will be the least extension of M to a c.t.m for ZFC containing G . The fact that $G \in M[G]$ will imply, by Lemma, that in most cases $M \neq M[G]$.

Let M be a c.t.m for ZFC, with \mathbb{P} a p.o. in M and G is \mathbb{P} – *generic* over M .

We shall show how to construct another c.t.m for ZFC, called $M[G]$, which will satisfy $M \subset M[G]$, $\alpha(M) = \alpha(M[G])$, and $G \in M[G]$.

$M[G]$ will be the least extension of M to a c.t.m for ZFC containing G . The fact that $G \in M[G]$ will imply, by Lemma, that in most cases $M \neq M[G]$.

Roughly, this will be the set of all sets which can be constructed from G by applying set-theoretic processes definable in M .

Let M be a c.t.m for ZFC, with \mathbb{P} a p.o. in M and G is \mathbb{P} – *generic* over M .

We shall show how to construct another c.t.m for ZFC, called $M[G]$, which will satisfy $M \subset M[G]$, $\alpha(M) = \alpha(M[G])$, and $G \in M[G]$.

$M[G]$ will be the least extension of M to a c.t.m for ZFC containing G . The fact that $G \in M[G]$ will imply, by Lemma, that in most cases $M \neq M[G]$.

Roughly, this will be the set of all sets which can be constructed from G by applying set-theoretic processes definable in M .

Each element of $M[G]$ will have a *name* in M , which tells how it has been constructed from G .

Let M be a c.t.m for ZFC, with \mathbb{P} a p.o. in M and G is \mathbb{P} – *generic* over M .

We shall show how to construct another c.t.m for ZFC, called $M[G]$, which will satisfy $M \subset M[G]$, $\alpha(M) = \alpha(M[G])$, and $G \in M[G]$.

$M[G]$ will be the least extension of M to a c.t.m for ZFC containing G . The fact that $G \in M[G]$ will imply, by Lemma, that in most cases $M \neq M[G]$.

Roughly, this will be the set of all sets which can be constructed from G by applying set-theoretic processes definable in M .

Each element of $M[G]$ will have a *name* in M , which tells how it has been constructed from G .

People living within M will be able to comprehend a name, τ , for an object in $M[G]$, but they will not in general be able to decide the object, τ_G , that τ names, since this will require a knowledge of G .

Definition

τ is a \mathbb{P} -name iff τ is a relation and

$$\forall \langle \sigma, p \rangle \in \tau [\sigma \text{ is a } \mathbb{P}\text{-name} \wedge p \in \mathbb{P}].$$

Definition

τ is a \mathbb{P} -name iff τ is a relation and

$$\forall \langle \sigma, p \rangle \in \tau [\sigma \text{ is a } \mathbb{P}\text{-name} \wedge p \in \mathbb{P}].$$

The collection of \mathbb{P} -name will be a proper class if $\mathbb{P} \neq 0$.

And this definition must be understood as a definition by transfinite recursion. Formally, one defines the characteristic function of the \mathbb{P} -name, $H(\mathbb{P}, \tau)$, by

$$H(\mathbb{P}, \tau) = 1 \text{ iff } \tau \text{ is a relation} \wedge \forall \langle \sigma, p \rangle \in \tau [H(\mathbb{P}, \sigma) = 1 \wedge p \in \mathbb{P}].$$

$$H(\mathbb{P}, \tau) = 0 \text{ otherwise.}$$

Definition

$V^{\mathbb{P}}$ is the class of \mathbb{P} -names. If M is a transitive model of ZFC and $\mathbb{P} \in M$, $M^{\mathbb{P}} = V^{\mathbb{P}} \cap M$. Or, by absoluteness,

$$M^{\mathbb{P}} = \{\tau \in M : (\tau \text{ is a } \mathbb{P}\text{-name})^M\}.$$

Definition

$V^{\mathbb{P}}$ is the class of \mathbb{P} -names. If M is a transitive model of ZFC and $\mathbb{P} \in M$, $M^{\mathbb{P}} = V^{\mathbb{P}} \cap M$. Or, by absoluteness,

$$M^{\mathbb{P}} = \{\tau \in M : (\tau \text{ is a } \mathbb{P}\text{-name})^M\}.$$

When forcing over M , use is made only of the \mathbb{P} -names in $M^{\mathbb{P}}$, which we may think of as being defined within M .

When $M \models \text{ZFC}$, τ is a \mathbb{P} -name $\Leftrightarrow (\tau \text{ is a } \mathbb{P}\text{-name})^M$.

Definition

$val(\tau, G) = \{val(\sigma, G) : \exists p \in G (\langle \sigma, p \rangle \in \tau)\}$. We also write τ_G for $val(\tau, G)$. And 0 is also a \mathbb{P} -name, we define that $0_G = 0$ for any G .

Definition

$val(\tau, G) = \{val(\sigma, G) : \exists p \in G (\langle \sigma, p \rangle \in \tau)\}$. We also write τ_G for $val(\tau, G)$. And 0 is also a \mathbb{P} -name, we define that $0_G = 0$ for any G .

Definition

If M is a transitive model of ZFC, $\mathbb{P} \in M$, and $G \subset \mathbb{P}$, then

$$M[G] = \{\tau_G : \tau \in M^{\mathbb{P}}\}.$$

Definition

$val(\tau, G) = \{val(\sigma, G) : \exists p \in G (\langle \sigma, p \rangle \in \tau)\}$. We also write τ_G for $val(\tau, G)$. And 0 is also a \mathbb{P} -name, we define that $0_G = 0$ for any G .

Definition

If M is a transitive model of ZFC, $\mathbb{P} \in M$, and $G \subset \mathbb{P}$, then

$$M[G] = \{\tau_G : \tau \in M^{\mathbb{P}}\}.$$

Lemma

If N is a transitive model of ZFC with $M \subset N$ and $G \in N$, then $M[G] \subset N$.

Thus, once we check that $M[G]$ is indeed a transitive extension of M containing G and satisfying ZFC, it will be the least such extension.

First, we want to check that $M \subseteq M[G]$:

Check Name

– \check{x} , $x \in M$ (check name)

$$\check{x} = \{ \langle \check{y}, \mathbb{1}_P \rangle \mid y \in x \}$$

e.g. $\check{0} = 0$, $\{\check{0}\} = \{ \langle 0, \mathbb{1}_P \rangle \}$

First, we want to check that $M \subseteq M[G]$:

Check Name

– \check{x} , $x \in M$ (check name)

$$\check{x} = \{ \langle \check{y}, \mathbb{1}_{\mathbb{P}} \rangle \mid y \in x \}$$

e.g. $\check{0} = 0$, $\{\check{0}\} = \{ \langle 0, \mathbb{1}_{\mathbb{P}} \rangle \}$

Lemma

If M is a transitive model of ZFC, \mathbb{P} is a p.o. in M , and G is a non-empty filter on \mathbb{P} , then:

- (a) $\forall x \in M [\check{x} \in M^{\mathbb{P}} \wedge \text{val}(\check{x}, G) = x]$.
- (b) $M \subset M[G]$.

Lemma

$$G \notin M \wedge G \in M[G].$$

Lemma

$$G \notin M \wedge G \in M[G].$$

Proof:

$(G \notin M)$ If not, $\mathbb{P} \setminus G \subset \mathbb{P}$. For \mathbb{P} , we know that $\forall p \in \mathbb{P} \exists q, r \in \mathbb{P}, q \perp r$. So $\mathbb{P} \setminus G$ is dense. Contradiction!

$(G \in M[G])$ we need to find a name that represents it.

$-\Gamma = \{ \langle \check{p}, p \rangle \mid p \in \mathbb{P} \}$, and $\Gamma_G = G$.

□

Lemma

$$G \notin M \wedge G \in M[G].$$

Proof:

$(G \notin M)$ If not, $\mathbb{P} \setminus G \subset \mathbb{P}$. For \mathbb{P} , we know that $\forall p \in \mathbb{P} \exists q, r \in \mathbb{P}, q \perp r$. So $\mathbb{P} \setminus G$ is dense. Contradiction!

$(G \in M[G])$ we need to find a name that represents it.

$-\Gamma = \{ \langle \check{p}, p \rangle \mid p \in \mathbb{P} \}$, and $\Gamma_G = G$.

□

Remark: $M[G]$ satisfy Axiom of Extension, Union, Pairing, Foundation.

Lemma

$M[G]$ is transitive.

Lemma

$M[G]$ is transitive.

Proof: $\tau_G \in M[G], \tau \in M^{\mathbb{P}}$
 $\tau_G = \{\sigma_G \mid \langle \sigma, p \rangle \in \tau, p \in G\}$
 So $\sigma_G \in M[G]$

□

Lemma

$M[G]$ is transitive.

Proof: $\tau_G \in M[G], \tau \in M^{\mathbb{P}}$
 $\tau_G = \{\sigma_G \mid \langle \sigma, p \rangle \in \tau, p \in G\}$
 So $\sigma_G \in M[G]$

□

Lemma

$Ord \cap M = Ord \cap M[G]$.

Lemma

$M[G]$ is transitive.

Proof: $\tau_G \in M[G], \tau \in M^{\mathbb{P}}$
 $\tau_G = \{\sigma_G \mid \langle \sigma, p \rangle \in \tau, p \in G\}$
 So $\sigma_G \in M[G]$

□

Lemma

$Ord \cap M = Ord \cap M[G]$.

Proof: $M \subseteq M[G]$, and they are transitive, so $Ord \cap M \subseteq Ord \cap M[G]$.
 For any set A , we define $rank(A) = \min\{\alpha \mid A \in V_\alpha\}$.
 $rk(A) = \sup\{rk(B) + 1 \mid B \in A\}$, and $\forall \alpha, rk(\alpha) = \alpha$.
 $Ord \cap M = \{rk(A) \mid A \in M\}$, $Ord \cap M[G] = \{rk(A) \mid A \in M[G]\}$.
 We want to prove that $\forall \tau, rk(\tau_G) \leq rk(\tau) \in Ord \cap M$, then we
 will know $\forall A \in M[G], rk(A) \in Ord \cap M$.

□

Now we should know that for a given φ if $M[G] \models \varphi$.

Now we should know that for a given φ if $M[G] \models \varphi$.
First, we define \Vdash :

$p \Vdash_{M, \mathbb{P}} \varphi$ iff for any (M, \mathbb{P}) – generic G , if $p \in G$, then $M[G] \models \varphi$.

That is a semantic definition, later we need to find a syntax definition. And then we prove that they are equivalent.

Let us consider first a specific example.

- M is a c.t.m for ZFC.
- \mathbb{P} is the set of finite partial functions from ω to 2 ordered by reverse inclusion.
- $\mathbb{1}_{\mathbb{P}}$ is the empty function.

And $\langle \mathbb{P}, \leq, \mathbb{1} \rangle \in M$, since its definition is absolute for transitive models of ZFC.

Let us consider first a specific example.

- M is a c.t.m for ZFC.
- \mathbb{P} is the set of finite partial functions from ω to 2 ordered by reverse inclusion.
- $\mathbb{1}_{\mathbb{P}}$ is the empty function.

And $\langle \mathbb{P}, \leq, \mathbb{1} \rangle \in M$, since its definition is absolute for transitive models of ZFC.

If G is a filter on \mathbb{P} , $f_G = \bigcup G$ is a function and $dom(f_G) = \omega$.

Let us consider first a specific example.

- M is a c.t.m for ZFC.
- \mathbb{P} is the set of finite partial functions from ω to 2 ordered by reverse inclusion.
- $\mathbb{1}_{\mathbb{P}}$ is the empty function.

And $\langle \mathbb{P}, \leq, \mathbb{1} \rangle \in M$, since its definition is absolute for transitive models of ZFC.

If G is a filter on \mathbb{P} , $f_G = \cup G$ is a function and $dom(f_G) = \omega$.

$G \in M[G]$ and $f_G = \cup G$, $f_G \in M[G]$ will follow immediately from the absoluteness of \cup for transitive models of ZF.

Let us consider first a specific example.

- M is a c.t.m for ZFC.
- \mathbb{P} is the set of finite partial functions from ω to 2 ordered by reverse inclusion.
- $\mathbb{1}_{\mathbb{P}}$ is the empty function.

And $\langle \mathbb{P}, \leq, \mathbb{1} \rangle \in M$, since its definition is absolute for transitive models of ZFC.

If G is a filter on \mathbb{P} , $f_G = \cup G$ is a function and $dom(f_G) = \omega$.

$G \in M[G]$ and $f_G = \cup G$, $f_G \in M[G]$ will follow immediately from the absoluteness of \cup for transitive models of ZF.

However, we may check $f_G \in M[G]$ directly, Let

$$\Phi = \{ \langle \langle n, m \rangle, p \rangle : p \in \mathbb{P} \wedge n \in dom(p) \wedge p(n) = m \}$$

$$\Phi_G = \{ \langle n, m \rangle : \exists p \in G (n \in dom(p) \wedge p(n) = m) \} = f_G.$$

Thus, $f_G \in M[G]$.

If G is \mathbb{P} – *generic* over M , then $G \notin M$.

For any $g : \omega \rightarrow 2$, if $g \in M$, $E = \{p : p \not\subseteq g\}$ is dense. So $G \cap E \neq \emptyset$ means that $f_G \neq g$. So $f_G \notin M$.

If G is \mathbb{P} – *generic* over M , then $G \notin M$.

For any $g: \omega \rightarrow 2$, if $g \in M$, $E = \{p: p \not\subseteq g\}$ is dense. So $G \cap E \neq \emptyset$ means that $f_G \neq g$. So $f_G \notin M$.

So in this example, we use forcing to find a new real number which is not in M .

If G is \mathbb{P} -generic over M , then $G \notin M$.

For any $g: \omega \rightarrow 2$, if $g \in M$, $E = \{p: p \not\subseteq g\}$ is dense. So $G \cap E \neq \emptyset$ means that $f_G \neq g$. So $f_G \notin M$.

So in this example, we use forcing to find a new real number which is not in M .

Also, we know that $\{ \langle 0, 0 \rangle \} \Vdash \dot{f}_G(0) = 0$ and $\{ \langle 0, 1 \rangle \} \Vdash \dot{f}_G(0) = 1$.

Now we define \Vdash^* .

$p \Vdash_{M, \mathbb{P}}^* \varphi(\tau_1, \tau_2, \dots, \tau_n)$, $\tau_1, \tau_2, \dots, \tau_n$ are names in $M^{\mathbb{P}}$, $\varphi(x_1, \dots, x_n)$ is a formula.

Now we define \Vdash^* .

$p \Vdash_{M, \mathbb{P}}^* \varphi(\tau_1, \tau_2, \dots, \tau_n)$, $\tau_1, \tau_2, \dots, \tau_n$ are names in $M^{\mathbb{P}}$, $\varphi(x_1, \dots, x_n)$ is a formula.

(Induction of the complexity of names and formula)

(a) $(\tau_1 = \tau_2)$, $\tau_1 = \{ \langle \pi_1, s_1 \rangle \mid \dots \}$, $\tau_2 = \{ \langle \pi_2, s_2 \rangle \mid \dots \}$.

(α) for all $\langle \pi_1, s_1 \rangle \in \tau_1$,

$\{q \leq p \mid q \leq s_1 \rightarrow \exists \langle \pi_2, s_2 \rangle \in \tau_2 (q \leq s_2 \wedge q \Vdash^* \pi_1 = \pi_2)\}$ is dense below p , and

(β) for all $\langle \pi_2, s_2 \rangle \in \tau_2$,

$\{q \leq p \mid q \leq s_2 \rightarrow \exists \langle \pi_1, s_1 \rangle \in \tau_1 (q \leq s_1 \wedge q \Vdash^* \pi_1 = \pi_2)\}$ is dense below p .

A is dense below p means that $\forall q \leq p \exists r \in A (r \leq q)$.

(b) $(\tau_1 \in \tau_2)$, $\tau_2 = \{ \langle \pi_2, s_2 \rangle \mid \dots \}$
 $\{ q \leq p \mid \exists \langle \pi_2, s_2 \rangle \in \tau_2 (q \leq s_2 \wedge q \Vdash^* \tau_1 = \pi_1) \}$ is dense below p .

(b) $(\tau_1 \in \tau_2), \tau_2 = \{ \langle \pi_2, s_2 \rangle \mid \dots \}$

$\{ q \leq p \mid \exists \langle \pi_2, s_2 \rangle \in \tau_2 (q \leq s_2 \wedge q \Vdash^* \tau_1 = \pi_1) \}$ is dense below p .

(c) $(\varphi \wedge \psi)$

$p \Vdash^* \varphi \wedge \psi$ iff $p \Vdash^* \varphi$ and $p \Vdash^* \psi$.

(b) $(\tau_1 \in \tau_2), \tau_2 = \{ \langle \pi_2, s_2 \rangle \mid \dots \}$

$\{ q \leq p \mid \exists \langle \pi_2, s_2 \rangle \in \tau_2 (q \leq s_2 \wedge q \Vdash^* \tau_1 = \pi_1) \}$ is dense below p .

(c) $(\varphi \wedge \psi)$

$p \Vdash^* \varphi \wedge \psi$ iff $p \Vdash^* \varphi$ and $p \Vdash^* \psi$.

(d) $(\neg \varphi)$

$p \Vdash^* \neg \varphi$ iff $\{ q \leq p \mid q \not\Vdash^* \varphi \}$ is dense below p .

(b) $(\tau_1 \in \tau_2)$, $\tau_2 = \{ \langle \pi_2, s_2 \rangle \mid \dots \}$
 $\{ q \leq p \mid \exists \langle \pi_2, s_2 \rangle \in \tau_2 (q \leq s_2 \wedge q \Vdash^* \tau_1 = \pi_1) \}$ is dense below p .

(c) $(\varphi \wedge \psi)$
 $p \Vdash^* \varphi \wedge \psi$ iff $p \Vdash^* \varphi$ and $p \Vdash^* \psi$.

(d) $(\neg \varphi)$
 $p \Vdash^* \neg \varphi$ iff $\{ q \leq p \mid q \not\Vdash^* \varphi \}$ is dense below p .

(e) $(\exists x \varphi)$
 $p \Vdash^* \exists x \varphi(x, \dots)$ iff $\{ r \leq p \mid \exists \sigma \in M^{\mathbb{P}} (r \Vdash^* \varphi(\sigma, \dots)) \}$ is dense below p .

Our goal: $p \Vdash_{M, \mathbb{P}}^* \varphi \Leftrightarrow p \Vdash_{M, \mathbb{P}} \varphi$

Our goal: $p \Vdash_{M, \mathbb{P}}^* \varphi \Leftrightarrow p \Vdash_{M, \mathbb{P}} \varphi$

Lemma

$$(p \Vdash \varphi) \wedge (q \leq p) \rightarrow q \Vdash \varphi$$

$$p \Vdash \varphi \text{ and } p \Vdash \psi \rightarrow p \Vdash \varphi \wedge \psi$$

Our goal: $p \Vdash_{M, \mathbb{P}}^* \varphi \Leftrightarrow p \Vdash_{M, \mathbb{P}} \varphi$

Lemma

$(p \Vdash \varphi) \wedge (q \leq p) \rightarrow q \Vdash \varphi$

$p \Vdash \varphi$ and $p \Vdash \psi \rightarrow p \Vdash \varphi \wedge \psi$

Lemma

TFAE:

- ① $p \Vdash^* \varphi(\vec{\tau})$.
- ② $\{q \leq p \mid q \Vdash^* \varphi(\vec{\tau})\}$ is dense below p .
- ③ $r \leq p, r \Vdash^* \varphi(\vec{\tau})$.

Proof:

(a)(b) By definition.

Proof:

(a)(b) By definition.

$$(1)_c \ p \Vdash^* \varphi \wedge \psi \\ \leftrightarrow \ p \Vdash^* \varphi \text{ and } p \Vdash^* \psi$$

Proof:

(a)(b) By definition.

$$(1)_c \ p \Vdash^* \varphi \wedge \psi$$

$$\leftrightarrow p \Vdash^* \varphi \text{ and } p \Vdash^* \psi$$

$$\leftrightarrow \forall r \leq p, r \Vdash^* \varphi \text{ and } r \Vdash^* \psi \text{ (By induction)}$$

Proof:

(a)(b) By definition.

$$(1)_c \ p \Vdash^* \varphi \wedge \psi$$

$$\leftrightarrow p \Vdash^* \varphi \text{ and } p \Vdash^* \psi$$

$$\leftrightarrow \forall r \leq p, r \Vdash^* \varphi \text{ and } r \Vdash^* \psi \text{ (By induction)}$$

$$\leftrightarrow \forall r \leq p, r \Vdash^* \varphi \wedge \psi \text{ (3)}_c$$

$$\rightarrow \{q \leq p \mid q \Vdash^* \varphi \wedge \psi\} \text{ is dense below } p. \text{ (2)}_c$$

$$\leftrightarrow \{q \leq p \mid q \Vdash^* \varphi \text{ and } q \Vdash^* \psi\} \text{ is dense below } p.$$

Proof:

(a)(b) By definition.

$$(1)_c \ p \Vdash^* \varphi \wedge \psi$$

$$\leftrightarrow p \Vdash^* \varphi \text{ and } p \Vdash^* \psi$$

$$\leftrightarrow \forall r \leq p, r \Vdash^* \varphi \text{ and } r \Vdash^* \psi \text{ (By induction)}$$

$$\leftrightarrow \forall r \leq p, r \Vdash^* \varphi \wedge \psi \text{ (3)}_c$$

$$\rightarrow \{q \leq p \mid q \Vdash^* \varphi \wedge \psi\} \text{ is dense below } p. \text{ (2)}_c$$

$$\leftrightarrow \{q \leq p \mid q \Vdash^* \varphi \text{ and } q \Vdash^* \psi\} \text{ is dense below } p.$$

$$\rightarrow \{q \leq p \mid q \Vdash^* \varphi\} \text{ is dense below } p \text{ and } \{q \leq p \mid q \Vdash^* \psi\} \text{ is dense}$$

below p.

$$\leftrightarrow p \Vdash^* \varphi \text{ and } p \Vdash^* \psi$$

$$\leftrightarrow (1)_c$$

(1) $p \Vdash^* \neg \varphi$
 $\leftrightarrow \{q \leq p \mid q \not\Vdash^* \varphi\}$ is dense below p .

(1) $p \Vdash^* \neg \varphi$

$\leftrightarrow \{q \leq p \mid q \not\Vdash^* \varphi\}$ is dense below p .

\leftrightarrow no $q \leq p$, $q \Vdash^* \varphi$

$(1)_d p \Vdash^* \neg \varphi$

$\Leftrightarrow \{q \leq p \mid q \not\Vdash^* \varphi\}$ is dense below p .

\Leftrightarrow no $q \leq p$, $q \Vdash^* \varphi$

$\rightarrow \forall r \leq p$, no $q \leq r$, $q \Vdash^* \varphi$

$\rightarrow (3)_d$

$\rightarrow (2)_d \{q \leq p \mid q \Vdash^* \neg \varphi\}$ is dense below p

$(1)_d p \Vdash^* \neg\varphi$

$\Leftrightarrow \{q \leq p \mid q \not\Vdash^* \varphi\}$ is dense below p .

\Leftrightarrow no $q \leq p$, $q \Vdash^* \varphi$

$\rightarrow \forall r \leq p$, no $q \leq r$, $q \Vdash^* \varphi$

$\rightarrow (3)_d$

$\rightarrow (2)_d \{q \leq p \mid q \Vdash^* \neg\varphi\}$ is dense below p

$\Leftrightarrow \{q \leq p \mid \text{no } r \leq q, r \Vdash^* \varphi\}$ is dense below p

$\Leftrightarrow (1)_d$

(1)_e $p \Vdash^* \exists x \varphi$

$\leftrightarrow \{r \leq p \mid \exists \sigma \in M^{\mathbb{P}}(r \Vdash^* \varphi)\}$ is dense below p

$(1)_e p \Vdash^* \exists x \varphi$

$\leftrightarrow \{r \leq p \mid \exists \sigma \in M^{\mathbb{P}}(r \Vdash^* \varphi)\}$ is dense below p

$\leftrightarrow \forall q \leq p, \{r \leq q \mid \exists \sigma \in M^{\mathbb{P}}(r \Vdash^* \varphi)\}$ is dense below q

$\leftrightarrow (3)_e$

$\rightarrow (2)_e \{r \leq p \mid r \Vdash^* \exists x \varphi\}$ is dense below p

$(1)_e \ p \Vdash^* \exists x \varphi$

$\leftrightarrow \{r \leq p \mid \exists \sigma \in M^{\mathbb{P}}(r \Vdash^* \varphi)\}$ is dense below p

$\leftrightarrow \forall q \leq p, \{r \leq q \mid \exists \sigma \in M^{\mathbb{P}}(r \Vdash^* \varphi)\}$ is dense below q

$\leftrightarrow (3)_e$

$\rightarrow (2)_e \ \{r \leq p \mid r \Vdash^* \exists x \varphi\}$ is dense below p

$\leftrightarrow \{r \leq p \mid \{q \leq r \mid \exists \sigma \in M^{\mathbb{P}}(q \Vdash^* \varphi)\}$ is dense below $r\}$ is dense below p

$\rightarrow (1)_e$

Theorem

- (1) $p \in G, p \Vdash^* \varphi(\tau_1, \dots, \tau_n)$
 $\Rightarrow M[G] \models \varphi(\tau_{1G}, \dots, \tau_{nG})$
- (2) $M[G] \models \varphi(\tau_{1G}, \dots, \tau_{nG})$
 $\Rightarrow \exists p \in G, p \Vdash^* \varphi(\tau_1, \dots, \tau_n)$

Theorem

- (1) $p \in G, p \Vdash^* \varphi(\tau_1, \dots, \tau_n)$
 $\Rightarrow M[G] \models \varphi(\tau_{1G}, \dots, \tau_{nG})$
- (2) $M[G] \models \varphi(\tau_{1G}, \dots, \tau_{nG})$
 $\Rightarrow \exists p \in G, p \Vdash^* \varphi(\tau_1, \dots, \tau_n)$

Corollary

- (1) $p \Vdash \varphi \Leftrightarrow p \Vdash^* \varphi$
- (2) $M[G] \models \varphi(\tau_{1G}, \dots, \tau_{nG}) \Leftrightarrow \exists p \in G (p \Vdash \varphi(\tau_1, \dots, \tau_n))$

Proof of Corollary:

- (1) (\Leftarrow) By the Theorem

(\Rightarrow) $p \Vdash \varphi$, we need to prove that $\{r \leq p \mid r \Vdash^* \varphi\}$ is dense below p .

If not, there is a G , s.t. $G \cap \{r \leq p \mid r \Vdash^* \varphi\} = \emptyset$. Then $M[G] \not\models \varphi$,

Contradiction!

- (1) \Rightarrow (2)

Part of the proof:

(1) $(\tau_1 = \tau_2)$ We assume $p \Vdash^* \tau_1 = \tau_2$ and $p \in G$. We must show $\tau_{1_G} = \tau_{2_G}$.

Part of the proof:

(1) ($\tau_1 = \tau_2$) We assume $p \Vdash^* \tau_1 = \tau_2$ and $p \in G$. We must show $\tau_{1_G} = \tau_{2_G}$.

We shall show $\tau_{1_G} \subset \tau_{2_G}$

Every element of τ_{1_G} is of the form π_{1_G} , where $\langle \pi_1, s_1 \rangle \in \tau_1$ for some $s_1 \in G$.

Part of the proof:

(1) $(\tau_1 = \tau_2)$ We assume $p \Vdash^* \tau_1 = \tau_2$ and $p \in G$. We must show $\tau_{1G} = \tau_{2G}$.

We shall show $\tau_{1G} \subset \tau_{2G}$

Every element of τ_{1G} is of the form π_{1G} , where $\langle \pi_1, s_1 \rangle \in \tau_1$ for some $s_1 \in G$.

Fix $r \in G$ with $r \leq p$ and $r \leq s_1$. Then $r \Vdash^* \tau_1 = \tau_2$, so there is a $q \in G$ such that $q \leq r$ and s.t. $q \leq s_1$ which implies

$$\exists \langle \pi_2, s_2 \rangle \in \tau_2 (q \leq s_2 \wedge q \Vdash^* \pi_1 = \pi_2). (*)$$

Part of the proof:

(1) $(\tau_1 = \tau_2)$ We assume $p \Vdash^* \tau_1 = \tau_2$ and $p \in G$. We must show $\tau_{1G} = \tau_{2G}$.

We shall show $\tau_{1G} \subset \tau_{2G}$

Every element of τ_{1G} is of the form π_{1G} , where $\langle \pi_1, s_1 \rangle \in \tau_1$ for some $s_1 \in G$.

Fix $r \in G$ with $r \leq p$ and $r \leq s_1$. Then $r \Vdash^* \tau_1 = \tau_2$, so there is a $q \in G$ such that $q \leq r$ and s.t. $q \leq s_1$ which implies

$$\exists \langle \pi_2, s_2 \rangle \in \tau_2 (q \leq s_2 \wedge q \Vdash^* \pi_1 = \pi_2). (*)$$

So fix $\langle \pi_2, s_2 \rangle$ as in (*), then $s_2 \in G$, so $\pi_{2G} \in \tau_{2G}$.

Also, by (1) for $\pi_1 = \pi_2$ (IH), $q \Vdash^* \pi_1 = \pi_2$ implies $\pi_{1G} = \pi_{2G}$, so $\pi_{1G} \in \tau_{2G}$.

The End